# Linear Decompositions for Multi-Valued Input Classification Functions 

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#### Abstract

In a multi-valued input classification function, each input combination represents properties of an object, while the output represents the class of the object. Each variable may have different radix. In most cases, the functions are partially defined. To represent multi-valued variables, both one-hot and minimum-length encoding are considered. Experimental results using University of California Irvine (UCI) benchmark functions show that the one-hot approach results in fewer variables than the minimum-length approach with linear decompositions.


## I. Introduction

We consider data sets with multi-valued inputs. Various methods exist to represent a multi-valued variable. To keep the original structure of the data, we use one-hot encoding. That is, to represent a $q$-valued variable, we use $q$ binary variables. Although this increases the total number of variables temporarily, the number of variables can be reduced by a linear decomposition. The original functions can be decomposed into two parts: a linear part, and a general part, as shown in Fig. 1.1. With this technique, many University of California Irvine (UCI) data sets for machine learning [21] were successfully decomposed. In contrast, conventional functional decompositions [1], [3], [9], [12] are more constrained and less likely to simplify. The rest of the paper is organized


Fig. 1.1. Linear Decomposition
as follows: Section II shows definitions and basic properties; Section III introduces linear decomposition; Section IV shows an illustrative example; Section V shows experimental results using UCI data sets; Section VI surveys related works; and Section VII concludes the paper.

## II. Definitions and Basic Properties

Definition 2.1: A multi-valued partially defined classification function is a mapping $f: D \rightarrow M$, where $D \subset Q^{N}$,

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TABLE 2.1
Registered Vector Table

| $x_{1}$ | $x_{2}$ | $x_{3}$ | $x_{4}$ | $x_{5}$ | $x_{6}$ | $f$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | 0 | 0 | 1 | 1 | 1 |
| 0 | 1 | 1 | 0 | 1 | 1 | 1 |
| 0 | 1 | 0 | 1 | 0 | 0 | 2 |
| 0 | 0 | 0 | 0 | 1 | 0 | 2 |
| 1 | 1 | 0 | 1 | 1 | 1 | 3 |
| 1 | 0 | 0 | 0 | 1 | 1 | 3 |

$Q=\{0,1,2, \ldots, q-1\}, M=\{1,2, \ldots, m\}$, and $N$ is the number of multi-valued variables. Each element of $D$ is called a registered vector. There are $k$ such vectors. $k$ is called the weight of the function. For each registered vector, assign an integer between 1 and $m$, where $2 \leq m \leq k$. A registered vector table shows the function value for each registered vector. A partially defined classification function $f$ produces the corresponding function value when the input vector matches a registered vector. Let $F_{i}$ be the set of registered vectors which map to $i \in M$. Then, $D=\bigcup_{i=1}^{m} F_{i}$. We assume that $F_{i} \neq \varnothing(i=1,2, \ldots, m)$.

Example 2.1: The registered vector table in Table 2.1 shows a classification function with $N=6, m=3$ and $k=6$. Here, the input variables are binary.

Definition 2.2: [4] For $F_{i}, F_{j} \subset Q^{N},(i, j=1,2, \ldots, m)$, and $i \neq j$, let $F_{i} \cap F_{j}=\varnothing$. An $m$-tuple $\left(F_{1}, F_{2}, \ldots, F_{m}\right)$ denotes a partially defined classification function when $\bigcup_{i=1}^{m} F_{i} \subset Q^{N}$. For a partially defined function $\left(F_{1}, F_{2}, \ldots, F_{m}\right)$, the function $\left(E_{1}, E_{2}, \ldots, E_{m}\right)$ that satisfies $F_{i} \subseteq E_{i} \subset Q^{N}$ is an extension of $\left(F_{1}, F_{2}, \ldots, F_{m}\right)$, where $E_{i} \cap E_{j}=\varnothing(i \neq j)$.

Definition 2.3: For a subset $U \subseteq Q^{N}$ and $S \subseteq$ $\{1,2, \ldots, N\}$, we denote by $\left.U\right|_{S}$ the projection of $U$ to $S$. In other words, $\left.U\right|_{S}=\left\{\left.\vec{a}\right|_{S}\right\}$, where $\left.\vec{a}\right|_{S}$ is obtained from $\vec{a} \in U$ by considering only those components $a_{j}$ with $j \in S$.

Example 2.2: Let $Q=\{0,1,2\}$ and $N=4$. Let $U=\{(1,2,0,1),(0,1,1,2),(2,0,1,2)\}$ and $S=\{2,3\}$. Then, we have the projection $\left.U\right|_{S}=\{(*, 2,0, *),(*, 1,1, *)$, $(*, 0,1, *)\}$.

Given a partially defined function, many extensions exist. In this paper, we seek an extension of $f$ that depends on the fewest variables.

Definition 2.4: Let $F_{i} \subset Q^{N}(i=1,2, \ldots, m)$. Given a partially defined function $\left(F_{1}, F_{2}, \ldots, F_{m}\right)$, and a subset $S \subseteq$ $\{1,2, \ldots N\}$, if $\left.\left.F_{i}\right|_{S} \cap F_{j}\right|_{S}=\varnothing, \quad(i \neq j)$ holds, then $S$

TABLE 2.2
Classification Function with Reduced Variables

| $x_{1}$ | $x_{2}$ | $x_{4}$ | $f$ | TAG | $y_{1}=x_{1}$ | $y_{2}=x_{2} \oplus x_{4}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | 0 | 1 | 1 | 1 | 1 |
| 0 | 1 | 0 | 1 | 2 | 0 | 1 |
| 0 | 1 | 1 | 2 | 3 | 0 | 0 |
| 0 | 0 | 0 | 2 | 4 | 0 | 0 |
| 1 | 1 | 1 | 3 | 5 | 1 | 0 |
| 1 | 0 | 0 | 3 | 6 | 1 | 0 |

is a support set. In such a case, $\left(\left.F_{1}\right|_{S},\left.F_{2}\right|_{S}, \ldots,\left.F_{m}\right|_{S}\right)$ is independent of the variable $x_{j}, j \in\{1,2, \ldots, n\}-S$, and the variable is redundant.

Example 2.3: Consider the function $\left(F_{1}, F_{2}, F_{3}\right)$ shown in Table 2.1. In this case, $S=\{1,2,4\}$ is a support set, since for

$$
\begin{aligned}
& \left.F_{1}\right|_{S}=\{(1,1, *, 0, *, *),(0,1, *, 0, *, *)\}, \\
& \left.F_{2}\right|_{S}=\{(0,1, *, 1, *, *),(0,0, *, 0, *, *)\}, \\
& \left.F_{3}\right|_{S}=\{(1,1, *, 1, *, *),(1,0, *, 0, *, *)\},
\end{aligned}
$$

$\left.\left.F_{i}\right|_{S} \cap F_{j}\right|_{S}=\varnothing$ holds for $i<j$. Thus, this function can be represented by three variables, as shown in Table 2.2.

Algorithms to represent a given function by using the minimum number of variables have been developed [11].

## III. Linear Decomposition

In this section, we assume that the input variables are binary. Let $n$ be the number of binary variables.

The number of variables of a partially defined classification function $f: D \rightarrow M$, where $D \subset B^{n}$ and $B=\{0,1\}$ often can be reduced by a linear decomposition [6], [8], [18]. In the linear decomposition shown in Fig. 1.1, $L$ realizes a linear function, while $G$ realizes a general function (in most cases, a non-linear function).

Definition 3.1: A compound variable has the form $y=$ $c_{1} x_{1} \oplus c_{2} x_{2} \oplus \cdots \oplus c_{n} x_{n}$, where $c_{i} \in\{0,1\}$. The compound degree of a variable $y$ is $\sum_{i=1}^{n} c_{i}$, where $\sum$ denotes ordinary integer addition, and $c_{i}$ is treated as an integer. A primitive variable is a variable whose compound degree is one.

When a partially defined function satisfies a certain condition, there exists a linear transformation that reduces the number of variables $p$.

Definition 3.2: In a partially defined function $\left(F_{1}, F_{2}, \ldots, F_{m}\right)$, let $\vec{a} \in F_{i}$, and $\vec{b} \in F_{j},(i \neq j)$. Then, the vector $\vec{d}=\vec{a} \oplus \vec{b}$ is a difference vector. The set of the difference vectors is denoted by $D_{f}$.

Lemma 3.1: [20] An $n$-variable partially defined classification function $f$ can be represented with $n-1$ compound variables if and only if there exists a non-zero vector $\vec{u}$ such that $\vec{u} \in B^{n}-D_{f}$, where $D_{f}$ is the set of difference vectors of $f$.

Theorem 3.1: [20] An $n$-variable partially defined classification function $f$ can be represented with at most $n-1$ compound variables if and only if the set of difference vectors $D_{f}$ contains less than $2^{n}-1$ distinct difference vectors.

Algorithm 3.1: (Reduction of Compound Variables [19])

TABLE 3.1
The set of difference vectors for the function in Table 2.2

| $x_{1}$ | $x_{2}$ | $x_{4}$ | Vector Pairs |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 1 | $(1,5),(2,3)$ |
| 0 | 1 | 0 | $(1,6),(2,4)$ |
| 1 | 0 | 0 | $(3,5), 4,6)$ |
| 1 | 0 | 1 | $(1,3),(2,5)$ |
| 1 | 1 | 0 | $(1,4),(2,6$ |
| 1 | 1 | 1 | $(3,6),(4,5)$ |

TABLE 3.2
CLASSIFICATION FUNCTION AFTER LINEAR TRANSFORMATION

| $y_{1}$ | $y_{2}$ | $f$ |
| :---: | :---: | :---: |
| 1 | 1 | 1 |
| 0 | 1 | 1 |
| 0 | 0 | 2 |
| 0 | 0 | 2 |
| 1 | 0 | 3 |
| 1 | 0 | 3 |

1) Derive the set of difference vectors $D_{f}$ of an $n$-variable function.
2) If $\left|D_{f}\right|=2^{n}-1$, then stop, since reduction is impossible.
3) Obtain a non-zero vector $\vec{d} \in B^{n}-D_{f}$ with minimum weight.
4) Remove one variable from $\vec{d}$, and apply the linear transformation to the function.
5) Let $n \leftarrow n-1$, and go to step 1 .

Example 3.1: Let us reduce the number of variables in the classification function shown in Table 2.2.

1) Table 3.1 shows the set of difference vectors. The rightmost column shows the pairs of registered vectors that produced the difference vectors. 12 difference vectors were generated, but only 6 vectors are distinct.
2) $\left|D_{f}\right|=6<2^{3}-1$.
3) In this case, the set of difference vectors does not contain $\vec{d}=(0,1,1)$. We select $\vec{d}=(0,1,1)$.
4) Perform the linear transformation that converts $\left(x_{2}, x_{4}\right)$ into $y_{2}=x_{2} \oplus x_{4}$. Thus, Table 2.2 is converted into Table 3.2, where $y_{1}=x_{1}$.
5) In this case, the function can be represented with two variables: $y_{1}=x_{1}$ and $y_{2}=x_{2} \oplus x_{4}$. The input columns in Table 3.2 show that $y_{1}$ and $y_{2}$ distinguish the three function values.

An efficient linear decomposition algorithm has been developed [19]. With this algorithm, functions with more than one thousand inputs have been successfully decomposed.

## IV. Illustrative Example

Our method to find multi-valued decompositions consists of two steps:

1) Represent the function using a binary encoding.
2) Perform a linear decomposition.

To show the method, consider the 2-variable 3-valued function, where the variables are 5 -valued, as shown in Table 4.1.

TABLE 4.1
5 -valued input 3-valued output function.

| $X_{1}$ | $X_{2}$ | $f$ |
| :---: | :---: | :---: |
| 2 | 4 | 0 |
| 0 | 3 | 0 |
| 0 | 1 | 0 |
| 3 | 1 | 1 |
| 2 | 3 | 1 |
| 2 | 2 | 1 |
| 4 | 0 | 2 |
| 3 | 3 | 2 |
| 1 | 4 | 2 |

## A. One-hot Encoding

In one-hot representations, $0,1,2,3,4$, and 5 become $(1,0,0,0,0),(0,1,0,0,0),(0,0,1,0,0),(0,0,0,1,0)$, and $(0,0,0,0,1)$, respectively. Similarly, the three output values for $f, 0,1$, and 2 , become $(1,0,0),(0,1,0)$, and $(0,0,1)$, respectively. With this substitution, Table 4.1 becomes Table 4.2.

TABLE 4.2
ONE-HOT ENCODED FUNCTION.

| $x_{0}$ | $x_{1}$ | $x_{2}$ | $x_{3}$ | $x_{4}$ | $y_{0}$ | $y_{1}$ | $y_{2}$ | $y_{3}$ | $y_{4}$ | $z_{0}$ | $z_{1}$ | $z_{2}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 0 | 0 |
| 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 1 | 0 | 0 |
| 1 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 0 |
| 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 0 |
| 0 | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 1 | 0 |
| 0 | 0 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |
| 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 1 |
| 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 |

The function in Table 4.2 is represented algebraically as

$$
Z=f(X, Y)
$$

where $X=\left(x_{0}, x_{1}, x_{2}, x_{3}, x_{4}\right), Y=\left(y_{0}, y_{1}, y_{2}, y_{3}, y_{4}\right)$, and $Z=\left(z_{0}, z_{1}, z_{2}\right)$. This function has a sum-of-products expression

$$
\begin{align*}
z_{0} & =x_{2} y_{4} \vee x_{0} y_{3} \vee x_{0} y_{1} \\
z_{1} & =x_{3} y_{1} \vee x_{2} y_{3} \vee x_{2} y_{2}  \tag{4.1}\\
z_{2} & =x_{4} y_{0} \vee x_{3} y_{3} \vee x_{1} y_{4}
\end{align*}
$$

Using Algorithm 3.1 yields the linear transformation:

$$
\begin{aligned}
w_{0} & =x_{3} \oplus y_{2} \oplus y_{3} \\
w_{1} & =x_{2} \oplus y_{1}
\end{aligned}
$$

In this case, $Z=\left(z_{0}, z_{1}, z_{2}\right)$ can be represented by only two variables, as shown in Table 4.3. Thus, the function can be represented by the following equations.

$$
\begin{align*}
z_{0} & =\bar{w}_{0} w_{1} \vee w_{0} \bar{w}_{1} \\
z_{1} & =w_{0} w_{1}  \tag{4.2}\\
z_{2} & =\bar{w}_{0} \bar{w}_{1}
\end{align*}
$$

Therefore, only two variables, $w_{0}$ and $w_{1}$, are needed between $L$ and $G$ in Fig. 1.1.

TABLE 4.3
One-hot encoded function: After linear transformation

| $w_{0}$ | $w_{1}$ | $z_{0}$ | $z_{1}$ | $z_{2}$ |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 1 | 0 | 0 |
| 1 | 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 0 | 0 |
| 1 | 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 1 | 0 |
| 0 | 0 | 0 | 0 | 1 |
| 0 | 0 | 0 | 0 | 1 |
| 0 | 0 | 0 | 0 | 1 |

## B. Minimum-length Encoding

In minimum-length encoding, a $q$-valued variable is represented by $\left\lceil\log _{2} q\right\rceil$ bits. We assume that the input part is represented by minimum-length encoding, while the output part is represented by one-hot encoding. Thus, Table 4.1 becomes Table 4.4. In this case, the number of variables can be

TABLE 4.4
Minimum-LENGTH ENCODED FUNCTION.

| $x_{2}$ | $x_{1}$ | $x_{0}$ | $y_{2}$ | $y_{1}$ | $y_{0}$ | $z_{0}$ | $z_{1}$ | $z_{2}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 0 | 1 | 1 | 1 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 | 1 | 0 | 0 |
| 0 | 1 | 1 | 0 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 0 |
| 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |
| 0 | 1 | 1 | 0 | 1 | 1 | 0 | 0 | 1 |
| 0 | 0 | 1 | 1 | 0 | 0 | 0 | 0 | 1 |

reduced to four variables: $x_{1}, x_{0}, y_{1}$, and $y_{0}$. Table 4.5 shows the function.

TABLE 4.5
Minimum-Length encoded function:After reduction of VARIABLES

| $x_{1}$ | $x_{0}$ | $y_{1}$ | $y_{0}$ | $z_{0}$ | $z_{1}$ | $z_{2}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 1 | 1 | 1 | 0 | 0 |
| 0 | 0 | 0 | 1 | 1 | 0 | 0 |
| 1 | 1 | 0 | 1 | 0 | 1 | 0 |
| 1 | 0 | 1 | 1 | 0 | 1 | 0 |
| 1 | 0 | 1 | 0 | 0 | 1 | 0 |
| 0 | 0 | 0 | 0 | 0 | 0 | 1 |
| 1 | 1 | 1 | 1 | 0 | 0 | 1 |
| 0 | 1 | 0 | 0 | 0 | 0 | 1 |

Algorithm 3.1 further reduced the number of variables to three as shown in Table 4.6, where $w=x_{0} \oplus y_{1}$.

TABLE 4.6
Minimum-LENGTH ENCODED FUNCTION:AFTER LINEAR TRANSFORMATION

| $x_{1}$ | $w$ | $y_{0}$ | $z_{0}$ | $z_{1}$ | $z_{2}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 0 | 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 | 0 | 0 |
| 0 | 0 | 1 | 1 | 0 | 0 |
| 1 | 1 | 1 | 0 | 1 | 0 |
| 1 | 1 | 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 0 | 1 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 |
| 1 | 0 | 1 | 0 | 0 | 1 |
| 0 | 1 | 0 | 0 | 0 | 1 |

## V. Experimental Results

Table 5.1 summarizes the benchmark functions used in the experiments. $N$ is the original number of variables; $m$ is the number of classes; $k$ is the number of instances, and $r$ is the radix. Original data were taken from the UCI (University of California, Irvine) machine learning repository [21].

Both one-hot encoding and minimum-length encoding were considered. First, original data were represented by one-hot encoding or minimum length encoding. Second, the linear decomposition algorithm in [19] was used to reduce the number of variables. Table 5.2 shows the experimental results. $n$ is the number of bits; $p_{0}$ is the number of bits after removing redundant bits; and $p$ is the number of bits after linear decomposition. The superior approach is indicated by a bold entry in Table 5.2. For these benchmark functions, one-hot encodings tend to produce smaller $p$. Over all 17 benchmark functions, the one-hot approach was superior to the minimum-length approach in 8 cases, while the minimum length approach was superior in one case. In 8 cases, both approaches produced the same number $p$ of variables.

The most time-consuming problem was Connect-4. It took 5709 seconds to reduce 126 variables into 22 variables using one-hot encoding. We used a computer with an Intel Core i7 7700 CPU, 4 cores, with 64 GB main memory, on Windows 10.

Brief descriptions of benchmark functions follow. Details are available in [21].

TABLE 5.1
Function Data [21]

| Data Name | $N$ | $m$ | $k$ | $r$ |
| :--- | ---: | ---: | ---: | ---: |
| Tic Tac Toe | 9 | 3 | 958 | 3 |
| Connect--4 | 42 | 3 | 67557 | 3 |
| Chess3196 | 36 | 2 | 3196 | $2,3,4$ |
| Chess28056 | 6 | 18 | 28056 | 4,8 |
| Poker | 10 | 10 | 25010 | 4,13 |
| Balance | 4 | 3 | 625 | 5 |
| MONK's | 6 | 2 | 432 | $2,3,4$ |
| Lymphography | 18 | 4 | 147 | $2,3,4,8$ |
| Breast Cancer | 9 | 2 | 677 | 10 |
| Hepatitis | 19 | 2 | 80 | $2 \sim 11$ |
| Nursery Schools | 8 | 5 | 12960 | $2,3,4,5$ |
| Car Evaluation | 6 | 4 | 1728 | 3,4 |
| Zoo | 16 | 7 | 101 | 2,6 |
| Vote | 16 | 2 | 435 | 3 |
| Promoter | 57 | 2 | 106 | 4 |
| Splice | 60 | 3 | 3174 | 4 |
| Mushrooms | 22 | 2 | 5644 | $2 \sim 12$ |

$N$ : Original number of variables $m$ : Number of classes $k$ : Number of instances
$r$ : Radix

## A. Tic-Tac-Toe

In tic-tac-toe, two players, X and Y , take turns marking the spaces in a $3 \times 3$ grid shown in Fig. 5.1. The player who first succeeds in placing three of their marks in a horizontal, vertical, or diagonal row is the winner [22]. This dataset encodes the complete set of possible board configurations at the end of a tic-tac-toe game, where X is assumed to have played first. The target concept is "win for $X$ ", i.e. true when X has one of 8 possible ways to create a "three-in-a-row". The number of instances is $k=958$ (legal tic-tac-toe endgame

TABLE 5.2
Comparison of Encodings.

|  | One-hot |  |  | Minimum-length |  |  |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: |
| Data Name | $n$ | $p_{0}$ | $p$ | $n$ | $p_{0}$ | $p$ |
| Tic Tac Toe | 27 | 9 | 9 | 18 | 12 | 9 |
| Connect-4 | 126 | 61 | 22 | 84 | 63 | 22 |
| Chess3196 | 75 | 30 | 15 | 38 | 30 | 15 |
| Chess28056 | 40 | 34 | 16 | 16 | 16 | 16 |
| Poker | 85 | 61 | $\mathbf{2 1}$ | 30 | 26 | 23 |
| Balance | 20 | 16 | $\mathbf{1 1}$ | 12 | 12 | 12 |
| MONK's | 17 | 5 | $\mathbf{3}$ | 10 | 8 | 4 |
| Lymphography | 59 | 11 | $\mathbf{8}$ | 29 | 11 | 9 |
| Breast Cancer | 90 | 13 | 10 | 36 | 11 | $\mathbf{9}$ |
| Hepatitis | 67 | 7 | $\mathbf{5}$ | 32 | 7 | 6 |
| Nursery Schools | 27 | 17 | $\mathbf{1 4}$ | 16 | 16 | 15 |
| Car Evaluation | 21 | 14 | 12 | 12 | 12 | 12 |
| Zoo | 36 | 6 | 5 | 19 | 7 | 5 |
| Vote | 48 | 11 | 8 | 32 | 11 | 8 |
| Promoter | 228 | 6 | $\mathbf{6}$ | 114 | 7 | 7 |
| Splice | 240 | 18 | $\mathbf{1 5}$ | 120 | 18 | 17 |
| Mushrooms | 118 | 5 | 4 | 53 | 7 | 4 |
| $\boldsymbol{n}$ Number of bits |  |  |  |  |  |  |
| $p: \quad$ Number of bits after removing blueundant bits |  |  |  |  |  |  |
| $p: \quad$ Number of bits after linear decomposition |  |  |  |  |  |  |

boards). The number of variables is $N=9$, and each variable takes $q=3$ values: (0) player X occupies, (1) player Y occupies, and (2) empty. The function takes $m=3$ classes ${ }^{1}$ : (0) player X wins, (1) player Y wins, and (2) draw. Among 958 instances, player X wins for 626 instances, player Y wins for 316 instances, and a draw occurs for 16 instances.

The function corresponds to a mapping: $D \rightarrow\{0,1,2\}$, where $D \subset\{0,1,2\}^{9}$, and $|D|=958$.

| $X_{1}$ | $X_{2}$ | $X_{3}$ |
| :--- | :--- | :--- |
| $X_{4}$ | $X_{5}$ | $X_{6}$ |
| $X_{7}$ | $X_{8}$ | $X_{9}$ |

Fig. 5.1. Board for tic-tac-toe
With one-hot encoding, this function was converted to a function with $n=N \times q=27$ binary variables. $n$ corresponds to the number of bits to represent the input part of the function. The number of primitive variables is reduced to $p_{0}=9$ variables.

## B. Connect-4

Connect-4 is a two-player board game. Two players, X and Y, drop disks into a vertically suspended $6 \times 7$ grid shown in Fig. 5.2. The disks fall straight down, occupying the lowest available space within the columns. The objective of the game is to be the first to form a a horizontal, vertical, or diagonal line of four of one's own disks.

This dataset contains all legal 8-ply (8 turn) positions in the game of Connect-4 in which neither player has won, and in which the next move is not forced. The number of variables is $N=6 \times 7=42$, since there are 6 rows and 7 columns in the grid. Each variable takes $q=3$ values, empty, X, and Y. The number of classes is $m=3:(0)$ player X wins, (1) player Y wins, and (2) draw. The number of instances is

[^0]

Fig. 5.2. Board for Connect-4
$k=67557$. Among these, player X wins for 44473 instances, player Y wins for 16635 instances, and a draw occurs for 6449 instances. The number of bits to represent the input part is $n=42 \times 3=126$. The number of primitive variables is reduced to $p_{0}=61$, and a linear decomposition reduced the number of compound variables to $p=22$,

## C. Chess 3196

Chess3196 is chess with $k=3196$ starting positions. The number of classes is $m=2$ : ( 0 ) white can win, and (1) white cannot win. The number of variables is $N=36: 34$ of them take two values; one of them takes three values; and one of them takes four values. Thus, the total number of binary variables is $n=(34 \times 2)+(1 \times 3)+(1 \times 4)=75$.

## D. Chess 28056

Chess28056 is chess with $k=28056$ starting positions. The number of classes is $m=18$. It shows the optimal depth-ofwin for White in 0 to 16 moves, otherwise draw. The number of variables is $N=6$ : two of which take four values; and four of which take eight values. Thus, the total number of binary variables is $n=(2 \times 4)+(4 \times 8)=40$.

## E. Poker

Each hand consists of five playing cards drawn from a standard deck of 52. Each card is described by two attributes (suit and rank). Thus, the number of variables is $N=10$ : five of which take four values (Hearts, Spades, Diamonds, Clubs); and five of which take 13 values (Ace, 2, 3, ... , Queen, King). Thus, the total number of binary variables is $n=(5 \times 4)+(5 \times 13)=85$. The class denotes the poker hand, and the number of classes is $m=10$ : (1) Nothing, (2) One pair, (3) Two pairs, (4) Three of a kind, (5) Straight, (6) Flush, (7) Full house, (8) Four of a kind, (9) Straight flush, and (10) Royal flush. The number of instances is $k=25010$.

## F. Balance Scale

This data models the operation of a balance scale. The number of classes is $m=3$ : (1) tip on the right, (2) balanced, and (3) tip on the left. $N=4$ variables are used: $X_{1}$ (the left weight), $X_{2}$ (the left distance), $X_{3}$ (the right weight), and $X_{4}$ (the right distance). Each variable takes one of $q=5$ values: $1,2,3,4,5$. The total number of binary variables is $n=4 \times 5=20$. The instances completely cover the attribute
space. Thus, the number of instances is $k=5^{4}=625$. The class is represented by

$$
f= \begin{cases}1 & \text { When }\left(X_{1} \times X_{2}\right)>\left(X_{3} \times X_{4}\right) \\ 2 & \text { When }\left(X_{1} \times X_{2}\right)=\left(X_{3} \times X_{4}\right) \\ 3 & \text { When }\left(X_{1} \times X_{2}\right)<\left(X_{3} \times X_{4}\right)\end{cases}
$$

## G. MONK's Problem

The MONK's problems were the basis of the first international comparison of learning algorithms. The number of classes is $m=2$. The number of variables is $N=6$ : two of them take two values; three of them take three values; and one of them takes four values. Thus, the total number of binary variables is $n=(2 \times 2)+(3 \times 3)+(1 \times 4)=17$. The instances completely cover the attribute space. Thus, the number of instances is $k=2^{2} \times 3^{3} \times 4=432$.

## H. Lymphography

This data came from the University Medical Center, Institute of Oncology, Ljubljana, Slovenia in Nov. 1988. The number of instances is $k=148$. The number of classes is $m=4$ : normal, metastases, malignant lymph, and fibrosis. The number of variables is $N=18$. Nine of them take two values; three of them take three values; four of them take four values; and one of them takes eight values. Thus, the total number of binary variables is $n=(9 \times 2)+(3 \times 3)+(4 \times 4)+(1 \times 8)=59$.

## I. Breast Cancer

This data came from the Clinical Sciences Center, University of Wisconsin. Features are computed from a digitized image of a fine needle aspiration (FNA) of a breast mass. The original dataset consists of 699 instances, including 16 incomplete instances, which are removed. We also removed the minimum number of instances so that the resulting dataset became consistent [17]. The number of remaining instances is $k=677$. The number of classes is $m=2$ : ( 0 ) malignant, and (1) benign. The number of variables is $N=9$, and each takes $q=10$ values. Thus, the total number of binary variables is $n=9 \times 10=90$.

## J. Hepatitis

This data came from Ljubljana, Slovenia in Nov. 1988. The number of instances is $k=80$. The number of classes is $m=2:(0)$ die and (1) live. The original dataset has 13 two-valued variables and 6 numerical variables. As for the numerical variables, 6 to 10 cut points are used to make multivalued variables. Thus, the resulting number of variables is $N=19: 13$ of which take two values; two of which take seven values; two of which take five values; one of which takes 6 values; and one of which takes 11 values. Thus, the total number of binary variables is $n=(13 \times 2)+(2 \times 7)+$ $(2 \times 5)+(1 \times 6)+(1 \times 11)=67$.

## K. Nursery Schools

A nursery school dataset was used to rank applications for nursery schools. It was used during 1980's when there was excessive enrollment in these schools in Ljubljana, Slovenia,
since rejected applications frequently needed an objective explanation. The number of classes is $m=5$ : not recommended, recommended, very recommended, priority, special priority. The number of variables is $N=8$ : one of which takes two values; four of which take three values; two of which take four values; and one of which takes five values. Thus, the total number of binary variables is $n=(1 \times 2)+(4 \times$ $3)+(2 \times 4)+(1 \times 5)=27$. The instances completely cover the attribute space. Thus, the number of instances is $k=2 \times 3^{4} \times 4^{2} \times 5=12960$.

## L. Car Evaluation

This data model for evaluating cars based on their price and technical characteristics. The number of classes is $m=$ 4: (0) unacceptable, (1) acceptable, (2) good, (3) very good. The number of variables is $N=6$ : three of which take three values; and three of which take four values. Thus, the total number of binary variables is $n=(3 \times 3)+(3 \times 4)=21$. The instances completely cover the attribute space. Thus, the number of instances is $k=3^{3} \times 4^{3}=1728$.

## M. Zoo

Given the properties of animals, this function shows the animal class. The number of classes is $m=7$ : (1) Mammals, (2) Birds, (3) Reptiles, (4) Fish, (5) Amphibians, (6) Insects, (7) Others. Properties of animals include: whether it drinks milk when it is a baby, whether it has feathers or not, number of legs etc. The number of variables is $N=16: 15$ of them take two values, and one of them takes 6 values. Thus, the total number of binary variables is $n=(15 \times 2)+(1 \times 6)=36$. The number of instances is $k=101$.

## N. Vote

This data is from the 1984 United States Congressional Voting Records. The number of instances, that is equal to the number of representatives, is $k=435$. The number of classes is $m=2$ : ( 0 ) democrat, and (1) republican. The number of variables is $N=16$; each takes three values: (0) yes, (1) no, (2) other. Thus, the total number of binary variables is $n=16 \times 3=48$.

## O. Promoter

This data is related to gene sequences (DNA), where the number of instances is $k=106$. The number of classes is $m=2$ : positive or negative. The number of variables is $N=$ 57, each takes four values: (1) A, (2) G, (3) T, and (4) C. Thus, the total number of binary variables is $n=57 \times 4=228$.

## P. Splice

This data is related to gene sequences (DNA). Originally, the data consists of 3191 instances. Removing 16 instances that had ambiguous inputs ( $\mathrm{D}, \mathrm{N}, \mathrm{S}, \mathrm{R}$ ), resulted in a data set with $k=3174$ instances. The number of classes is $m=3$ : (0) donor, (1) acceptor, (2) neither. The number of variables is $N=60$; each takes four values: (1) A, (2) G, (3) T, and (4) C. Thus, the total number of binary variables is $n=60 \times 4=240$.

## Q. Mushroom

This dataset includes descriptions of hypothetical samples corresponding to 23 species of gilled mushrooms in the Agaricus and Lepiota Family. Each species is identified as definitely edible, definitely poisonous, or of unknown edibility and not recommended. This latter class was combined with the poisonous class. The original data contains 8124 mushrooms; among them, 2480 are incomplete, which are removed. Thus, the number of instances is $k=5644$. The number of classes is $m=2$ : (1) edible, and (2) not edible. The number of variables is $N=22$. Among them, four variables take two values; two take three values; five take four values; three take six values; one takes seven values; one takes eight values; four take nine values; one takes 10 values; and one takes 12 values. Thus, the total number of binary variables is $n=(4 \times 2)+(2 \times 3)+(5 \times$ $4)+(3 \times 6)+(1 \times 7)+(1 \times 8)+(4 \times 9)+(1 \times 10)+(1 \times 12)=125$. Since, we removed incomplete data, the number of binary variables was reduced to 118 .

## VI. Related Works

Decompositions of multi-valued input functions In [23], many multi-valued benchmark functions were decomposed to show the usefulness for machine learning. However, they did not show the numbers of bits after decompositions in the table. But, they did show that the numbers of reduced attributes (multi-valued variables) for MONK's, vote, splice, and mushrooms are $3,10,12$, and 5 , respectively.

In [5], bi-decompositions of multi-valued functions were considered. In [16], an efficient optimum functional decompositions algorithm was developed for index generation functions. In [14], [15], linear decompositions of multi-valued input index generation functions were considered, where $\bmod (p)$ additions were used in the linear transformations.

Support reduction of multi-valued input functions In [7], BDDs were used to reduce the support of binary encoded multi-valued functions. This is the most relevant work to our research. They did not use the UCI data set, but used the POLO benchmark set, which is not available now. So, we cannot check if the data are the same. They used minimum-length encoding to represent multi-valued input and output values. On the other hand, we used one-hot encoding to represent the output values. So, direct comparison is difficult. As for zoo and mushrooms, the results are the same. As for tic-tac-toe, they obtained an $p=8$-bit solution.

In [2], one-hot encodings were used to find essential variables. In [17], variable minimization of multi-valued input binary functions (i.e., $m=2$ ) was considered.

## VII. CONCLUSION AND COMMENTS

This paper shows a method to decompose multi-valued input classification functions. First, original functions were converted into partially defined functions. Then, linear decompositions were used to reduce the number of variables. With this strategy, functions with many variables were decomposed successfully. Both one-hot encoding and minimum-length encoding were considered. Over all 17 benchmark functions,
the one-hot approach was superior to the minimum-length approach in 8 cases, while the minimum length approach was superior in one case. In 8 cases, both approaches produced the same number $p$ of variables. The superior approach is indicated by a bold entry in Table 5.2.

In a one-hot encoding, a function has more variables in the initial representation. This means that more choices exist for a linear transformation to reduce the variables. On the other hand, in a minimum-length encoding, fewer choices exist, and the number of applicable linear transformations is smaller. Thus, the one-hot encodings tend to result in fewer variables. This result is consistent with the experimental results in index generation functions [18]. That is, when the number of registered vectors are the same, randomly generated functions with more variables in the original representations tend to result in a fewer variables by linear transformations.

This method is promising for data mining.

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[^0]:    ${ }^{1}$ The function does not show how to win. Instead, it represents a win for X , a win for Y or, a win for neither, assuming each player plays optimally.

